

CSCI 210: Computer Architecture

Lecture 24: Datapath

Stephen Checkoway

Oberlin College

Apr. 22, 2022

Slides from Cynthia Taylor

Announcements

- Problem Set 7 due today
- Lab 6 due Sunday
- Office Hours today 13:30–14:30 pm

Amdahl's Law

$$\text{Execution time after improvement} = \frac{\text{Execution Time Affected}}{\text{Amount of Improvement}} + \text{Execution Time Unaffected}$$

Amdahl's law example

- A program originally takes **30 seconds** to run
- **One third** of the execution time comes from a single loop
- With some clever programming, you speed the loop up such that it runs **twice as fast**
- How long does the improved program take to run?

$$\text{Execution time after improvement} = \frac{\text{Execution Time Affected}}{\text{Amount of Improvement}} + \text{Execution Time Unaffected}$$

What was the overall speedup due to the improved loop?

- The original program took 30 s
- The new program took 25 s
- Speedup = Original Time / New Time
- Speedup = 30 s / 25 s = 1.2

Amdahl's Law and Parallelism

Our program is **90% parallelizable** (segment of code executable in parallel on multiple processor cores) and runs in **100 seconds** with a single core. What is the execution time if you use **4 cores** (assume no overhead for parallelization)?

	Execution Time
A	25 seconds
B	32.5 seconds
C	50 seconds
D	92.5 seconds
E	None of the above

$$\text{Execution time after improvement} = \frac{\text{Execution Time Affected}}{\text{Amount of Improvement}} + \text{Execution Time Unaffected}$$

Amdahl's Law

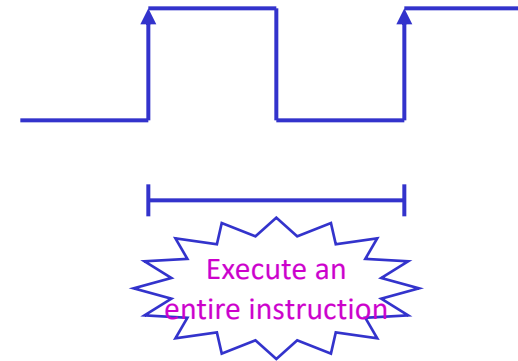
So what does Amdahl's Law *mean* at a high level?

Selection	"BEST" message from Amdahl's Law
A	Parallel programming is critical for improving performance
B	Improving serial code execution is ultimately the most important goal.
C	Performance is strictly tied to the ability to determine which percentage of code is parallelizable.
D	The impact of a performance improvement is limited by the percent of execution time affected by the improvement
E	None of the above

Performance Questions?

The Big Picture: The Performance Perspective

- Processor design (datapath and control) will determine:
 - Clock cycle time
 - Clock cycles per instruction
- Starting today:
 - Single cycle processor:
 - Advantage: One clock cycle per instruction
 - Disadvantage: long cycle time
- $ET = Insts * CPI * Cycle Time$



The Processor: Datapath & Control

- We're ready to look at an implementation of MIPS simplified to contain only:
 - memory-reference instructions: `lw, sw`
 - arithmetic-logical instructions: `add, sub, and, or, slt`
 - control flow instructions: `beq`

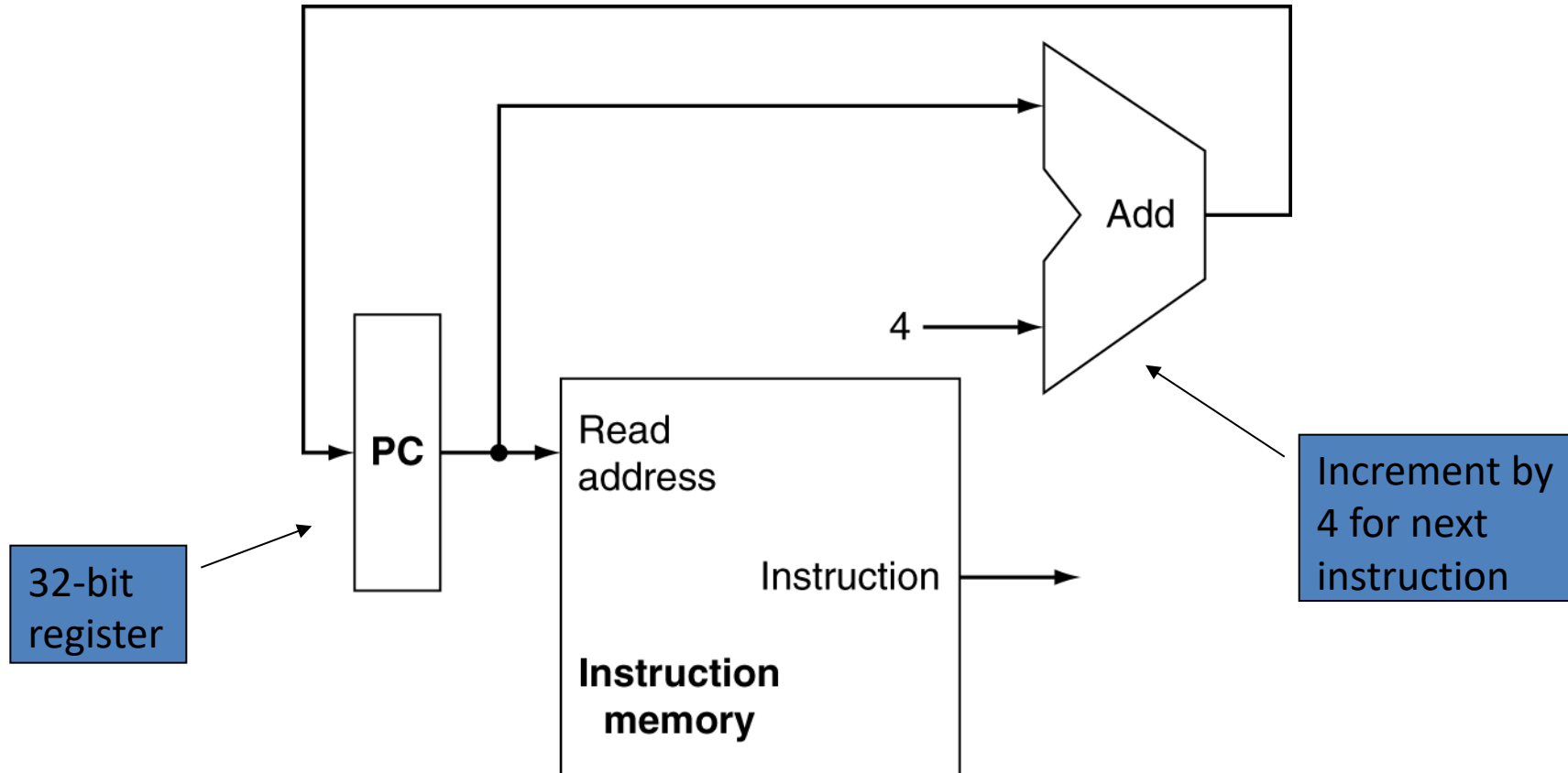
Generic implementation

- Fetch
 - Use the program counter (PC) to supply instruction address
 - Get the instruction from memory
 - Update the program counter to the next instruction
- Decode instruction
 - Read registers
 - Use the instruction to decide exactly what to do
- Execute
 - Perform necessary data manipulation
 - Write to registers

To fetch an instruction, what hardware do we need?

- Fetch
 - Use the program counter (PC) to supply instruction address
 - Get the instruction from memory
 - Update the program counter to the next instruction
- A. Register(s), Memory
 - B. Register(s), Adder, Memory
 - C. Register(s), ALU, Memory
 - D. More than this

Instruction Fetch



Generic implementation

- Fetch
 - Use the program counter (PC) to supply instruction address
 - Get the instruction from memory
 - Update the program counter to the next instruction
- **Decode instruction**
 - Read registers
 - Use the instruction to decide exactly what to do
- Execute
 - Perform necessary data manipulation
 - Write to registers

Which of these describes the interface for our register file?

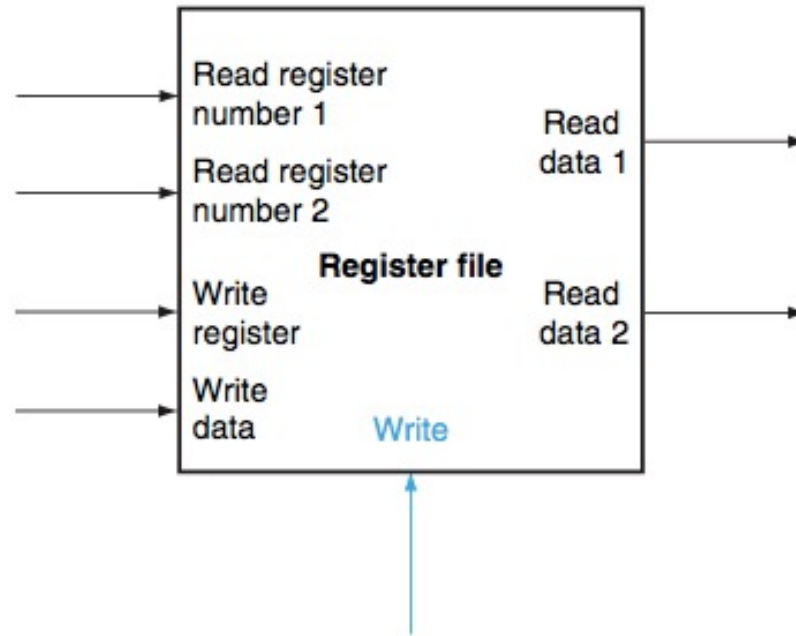
Think about what we will read in and out in different instructions i.e.,

```
add $t0, $t0, $t1 vs.
```

```
sw $t0, 16($s3)
```

- A. Two 32-bit data outputs, three 5-bit select inputs, 1-bit control input
- B. Two 32-bit data outputs, three 32-bit select inputs, 1-bit control input
- C. Two 32-bit data outputs, three 5-bit select inputs, 1 32-bit data input, 1-bit control input
- D. Two 32-bit data outputs, 2 32-bit select inputs, 1 5-bit data input, 1-bit control input
- E. None of the above

Register File



Reading

- Next lecture: Control Path
 - Section 5.3